

Improving Relational Consistency Algorithms Using Dynamic Relation Partitioning

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Constraint Systems Laboratory

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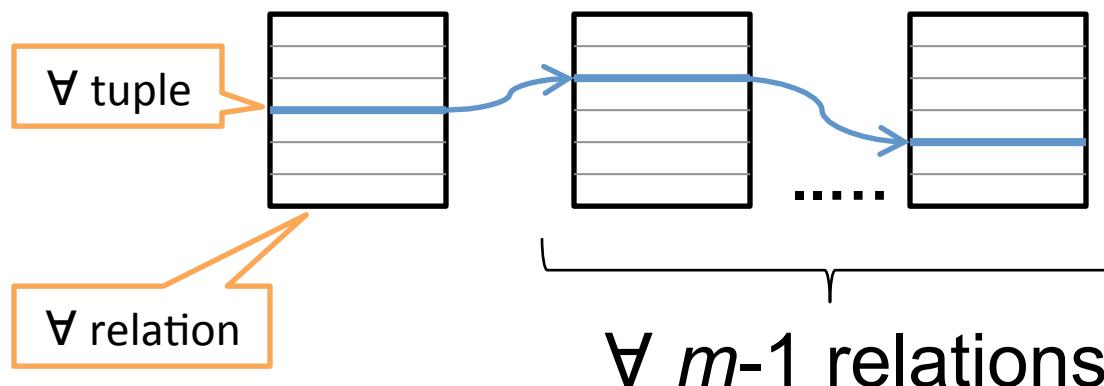


Outline

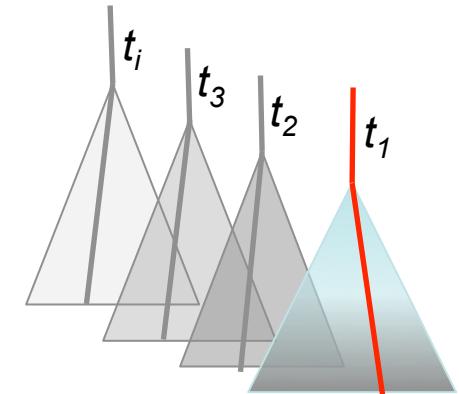
- Introduction
 - $R(*,m)C$ property and its algorithm PERTUPLE
- Partitioning a relation
 - Coarse, fine, intermediate blocks
- Improve PERTUPLE using partitions
 - PERTUPLE → PERFB
- Experimental results
- Conclusion

$R(*,m)C$ (a.k.a. m -wise consistency)

- A CSP is $R(*,m)C$ iff
 - Every tuple in a relation can be extended
 - to the variables in the scope of any ($m-1$) other relations
 - in an assignment satisfying all m relations simultaneously
- $R(*,m)C \equiv$ Every set of m relations is minimal



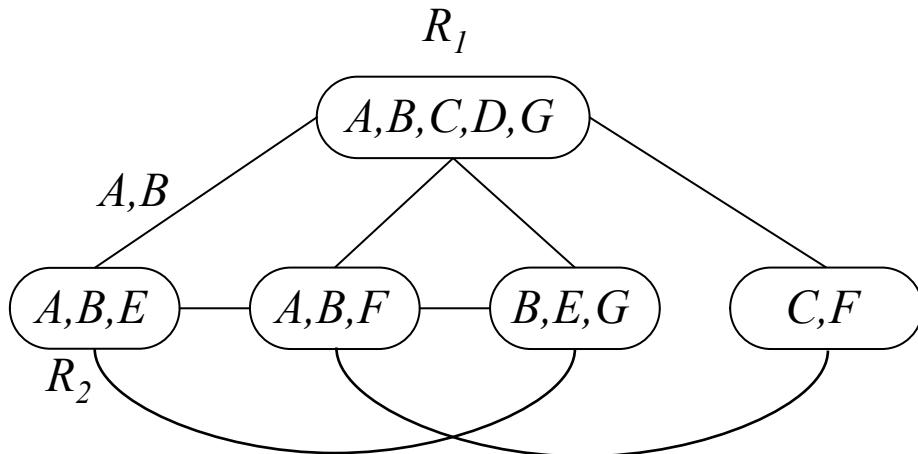
- PERTUPLE enforces $R(*,m)C$
- Store all connected combinations of m relations
- For each relation in a combination
 - For each tuple in the relation
 - **SEARCHSUPPORT**: Conduct backtrack search with FC over the dual CSP induced by the m relations
 - Remove the tuple if no solution is found
- Update propagation queue



Piecewise Functional Constraints

- PW-AC enforces $R(*,2)C$

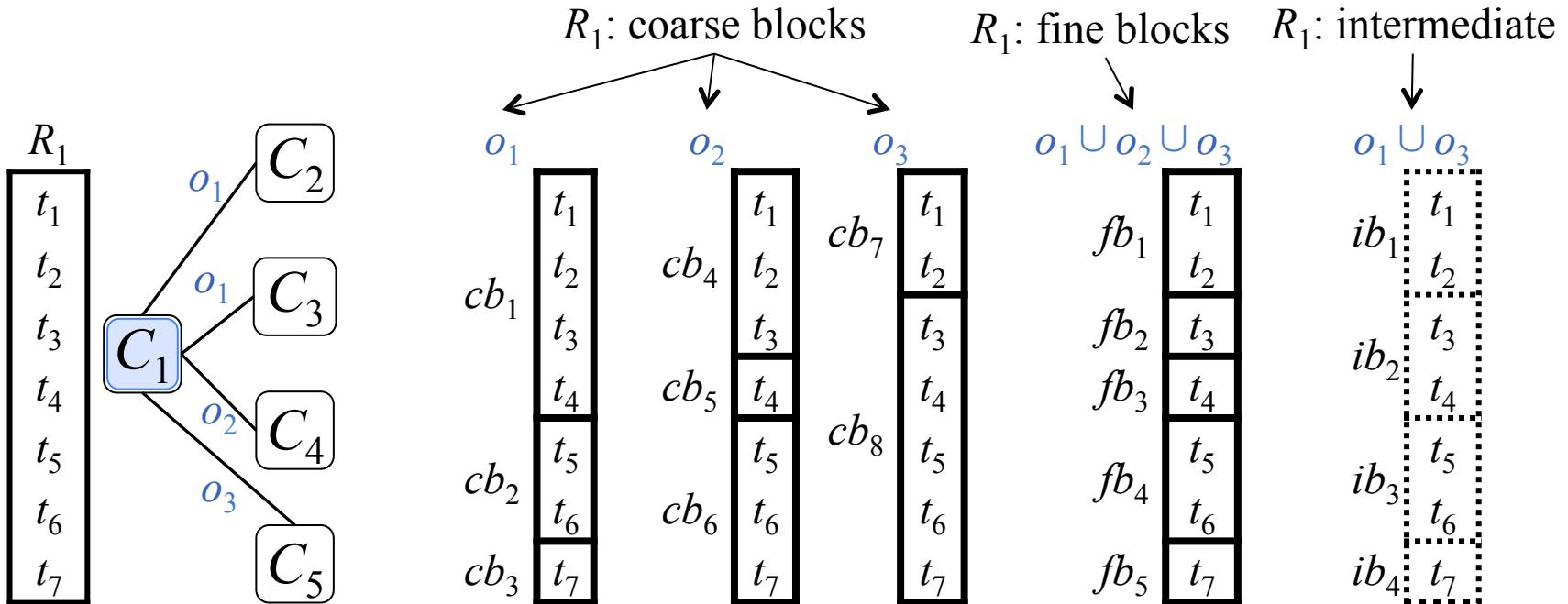
[Samaras & Stergiou JAIR 05]



	A	B	C	D	G
t_1	0	0	0	0	0
t_2	0	0	0	1	0
t_3	0	0	1	0	0
t_4	0	0	1	1	1
t_5	0	1	1	0	1
t_6	0	1	1	1	1
t_7	1	1	1	1	1

	A	B	E
t_1	0	0	0
t_2	0	0	1
t_3	0	1	0
t_4	0	1	1
t_5	1	0	0
t_6	1	0	1

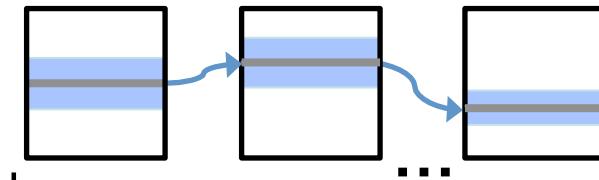
Types of Partitions



- For each relation, we store
 - A single partition of fine blocks
 - As many partitions of coarse blocks as shared subscopes

PERTUPLE → PERFB

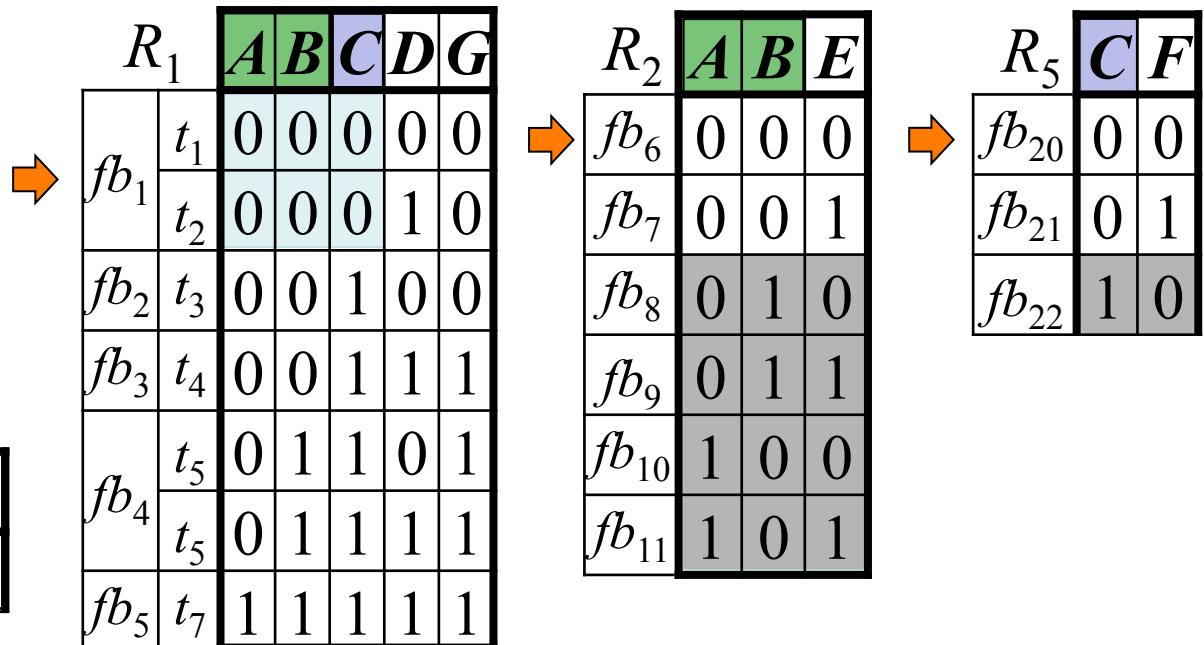
- Tuple → Fine Block (FB)
- General mechanism is identical
 - Combinations, queuing, and propagation
- **SEARCHSUPPORT**
 - Search enumerates fine blocks, not tuples
 - Forward checking operates on coarse blocks
- Calls to **SEARCHSUPPORT** reduced
 - Skips tuples in the same fine blocks
 - Skips fine blocks in the same intermediate block
 - Using two local data structures
- Details in paper



Enforcing R(*,m)C with PERFB

Vars in subscopes

Rel	Variables
.	



Equiv FBs for R₁

A	B	C	Consistent
0	0	0	True

Enforcing R(*,m)C with PERFB

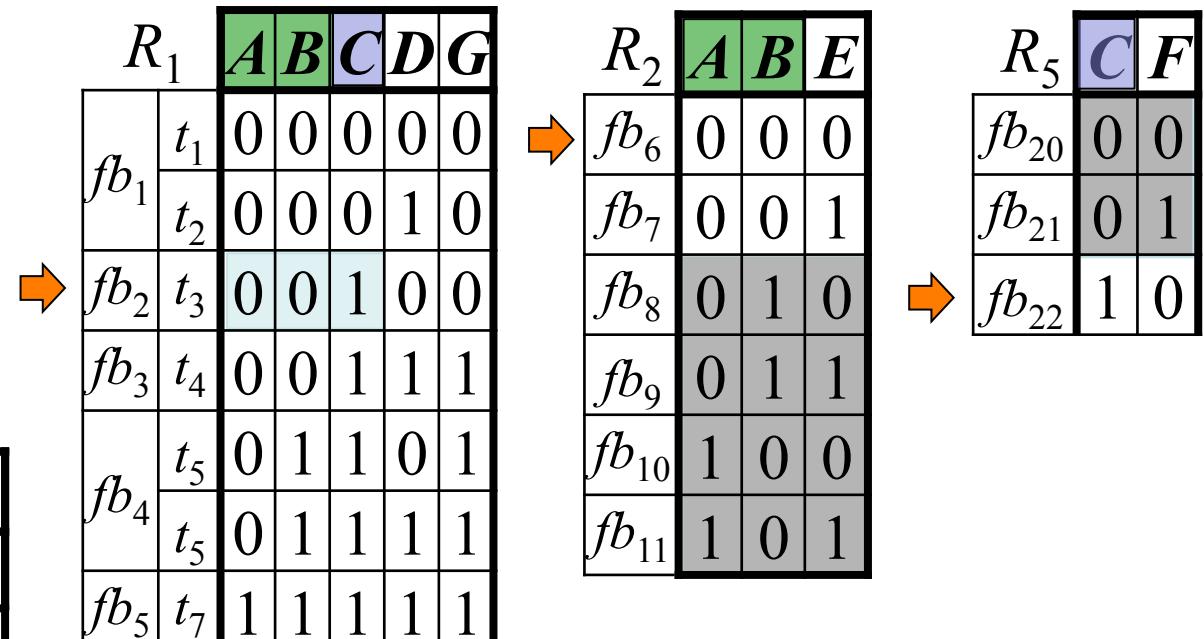
Vars in subscopes

Rel	Variables
R_1	ABC

?

Equiv FBs for R_1

A	B	C	Consistent
0	0	0	True
0	0	1	True



Enforcing R(*,m)C with PERFB

Vars in subscopes

Rel	Variables
R_1	ABC

Yes!

Equiv FBs for R_1

A	B	C	Consistent
0	0	0	True
0	0	1	True



R_1		A	B	C	D	G
fb_1	t_1	0	0	0	0	0
fb_1	t_2	0	0	0	1	0
fb_2	t_3	0	0	1	0	0
fb_3	t_4	0	0	1	1	1
fb_4	t_5	0	1	1	0	1
fb_4	t_5	0	1	1	1	1
fb_5	t_7	1	1	1	1	1

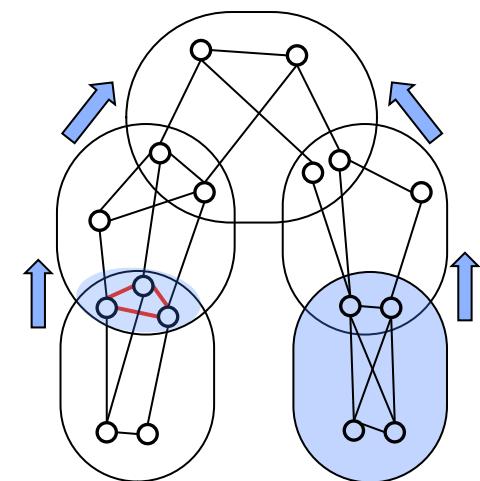
R_2		A	B	E
fb_6		0	0	0
fb_7		0	0	1
fb_8		0	1	0
fb_9		0	1	1
fb_{10}		1	0	0
fb_{11}		1	0	1

R_5	C	F
fb_{20}	0	0
fb_{21}	0	1
fb_{22}	1	0

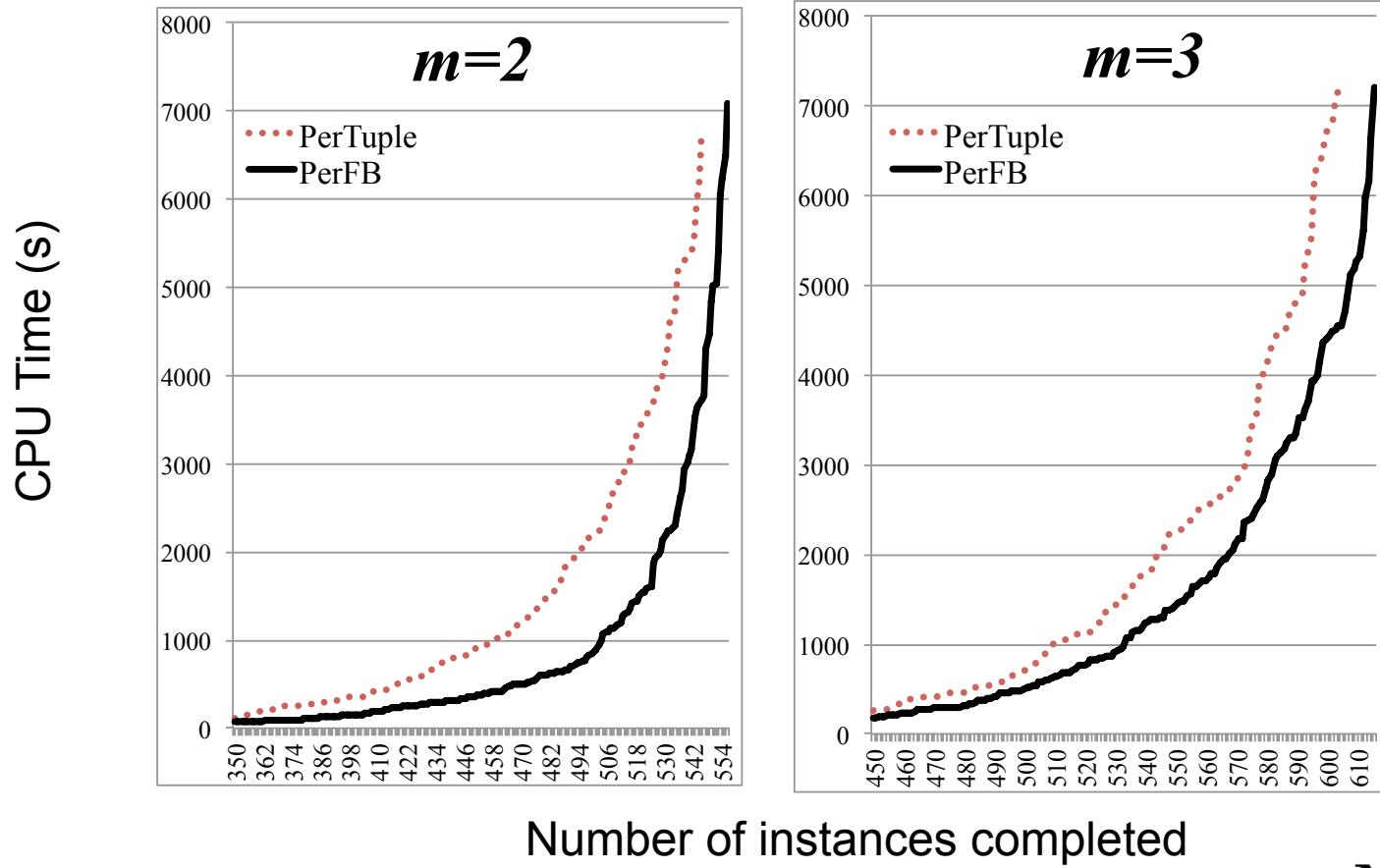
Experimental Setup

- Backtrack search, find first solution, dom/deg
- Maintaining $R(*,m)C$
 - Use a tree-decomposition of the CSP
 - Enforce consistency on individual clusters
 - Add projection of constraints to separators to bolster propagation between adjacent clusters
 - Use a minimal dual graph to reduce number of combinations of m constraints
 - Better performance than GAC, maxRPWC
- $m = 2, 3, 4, |\psi(\text{cl})|$
- XCSP benchmark of CP Solver Competition

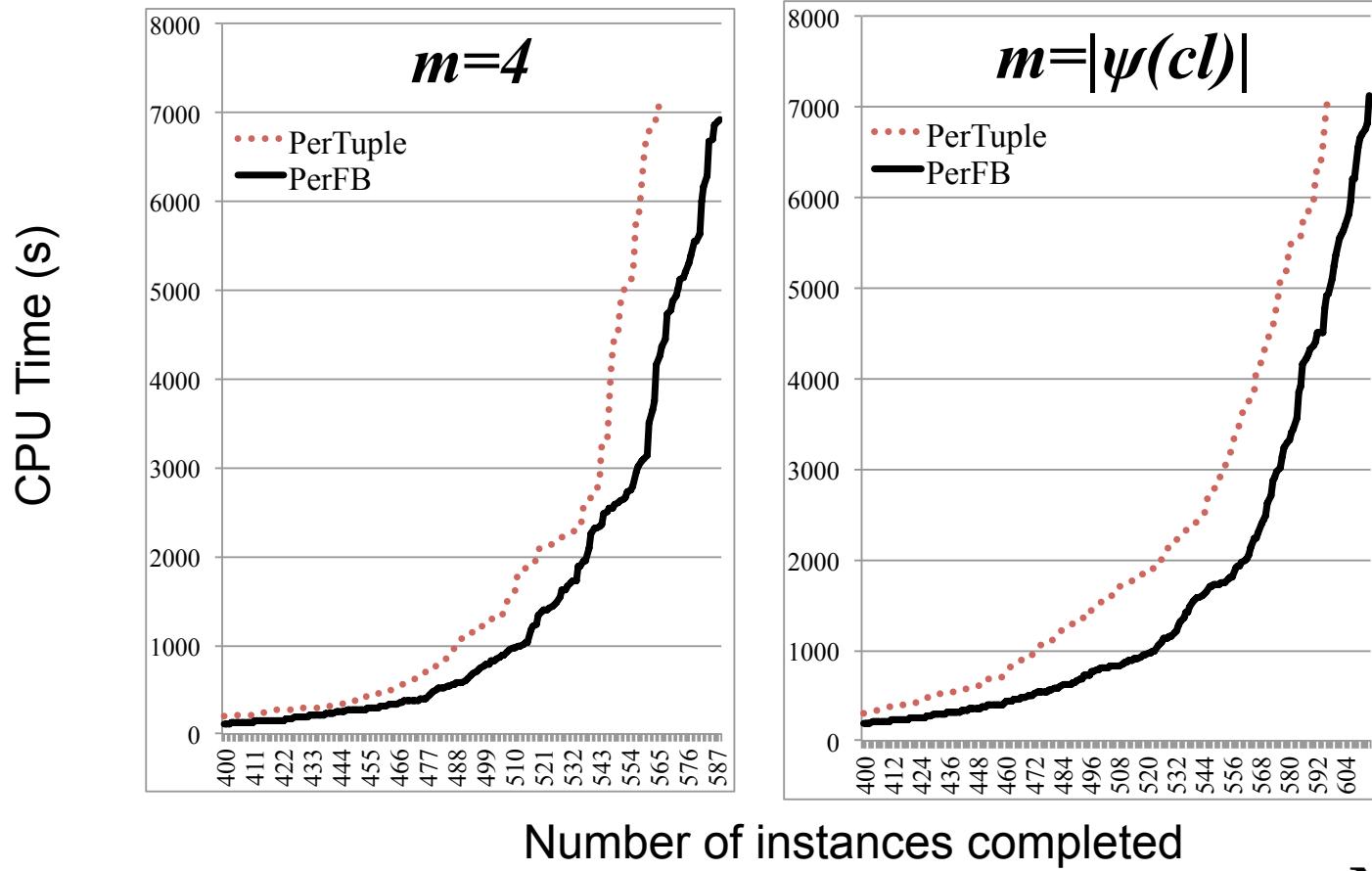
[Karakashian+ AAAI 13]



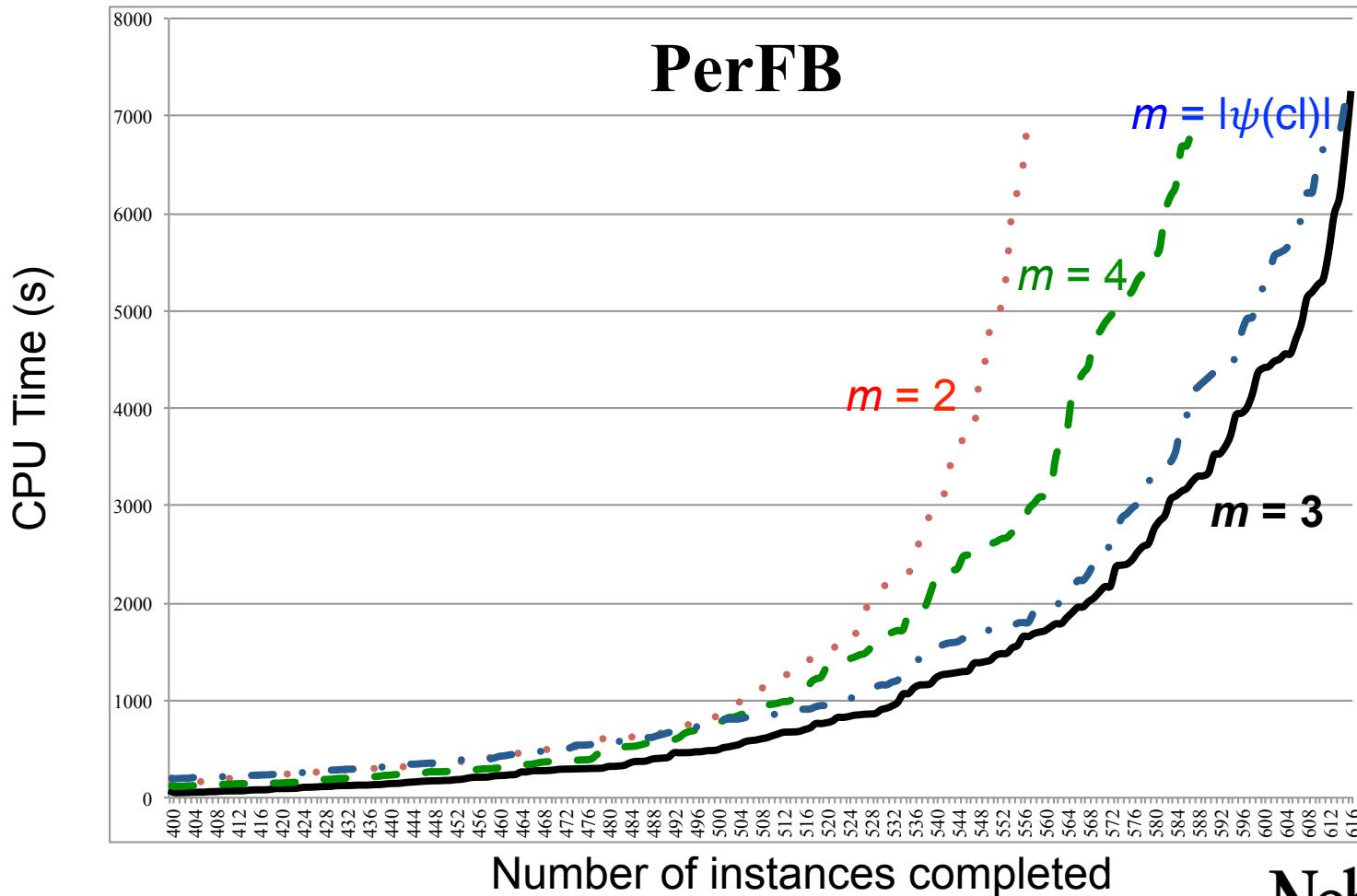
PERFB vs. PERTUPLE: completed instances



PERFB vs. PERTUPLE: completed instances



PERFB $m = 2,3,4, |\psi(\text{cl})|$



Detailed Results

Number of instances tested: 853

	$m = 2$	$m = 3$	$m = 4$	$m = \psi(\text{cl}) $
	PERTUPLE	PERFB	PERTUPLE	PERFB
#Completed	546	557	604	616
... only by	5	16	1	13
... by both	541		603	
Avg. CPU (sec)	538	227	521	362
SEARCHSUPPORT calls (10^9)	86.4	0.0	88.1	26.1
Call ratio	--		3.37	2.69
				3.06

Conclusions

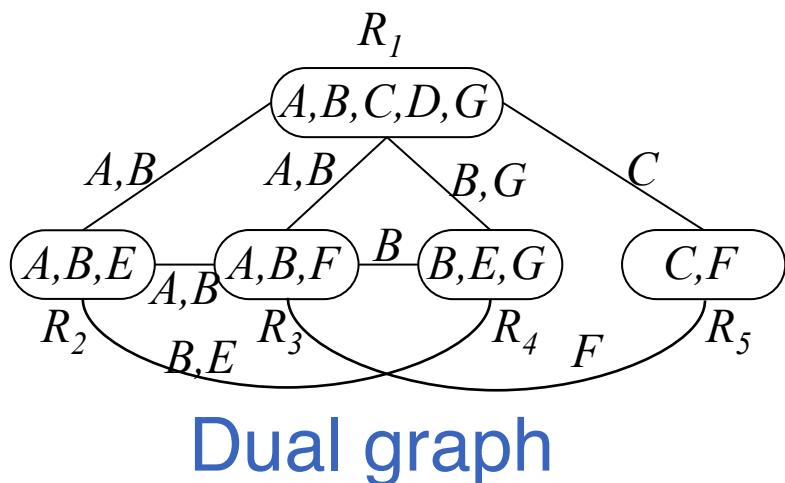
- Contributions
 - Designed PERFB
 - to replace PERTUPLE of [Karakashian+ AAAI 10]
 - by extending the work of [Samaras & Stergiou JAIR 05]
 - Empirically showed benefits of our approach
- Future work
 - Extend our approach to our other algorithm for enforcing $R(*,m)C$ [Karakashian PhD 13]

Thank you for listening

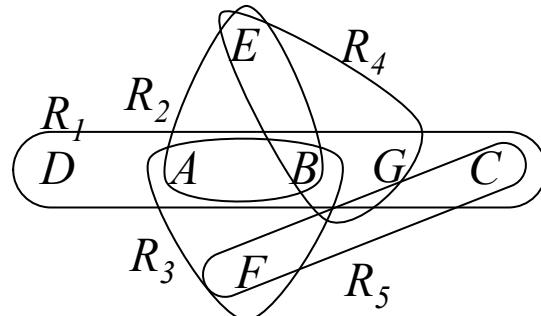
Wake up the Chair!

CSP: Graphical Representations

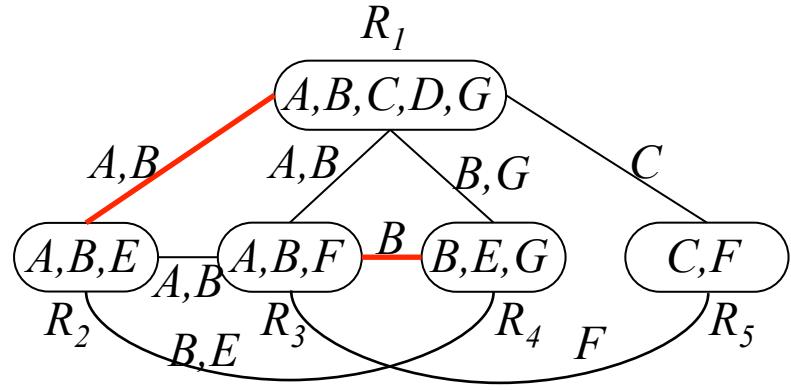
- Hyper graph
- Dual graph
- Minimal dual graph



Dual graph



Hyper graph



Minimal dual graph

[Janssen+, 1989]

Block Statistics

Benchmark	Absolute Max	Averages		
		Min	Max	Mean
geom	17	1.0	1.2	1.0
graphColoring-hos	3	1.0	2.0	1.0
graphColoring-sgb-book	12	1.0	7.7	1.1
hanoi	2	1.0	2.0	1.0
modifiedRenault	260	1.0	25.6	1.0
rand-10-20-10	2	1.0	1.3	1.0
renault	4	1.0	4.0	1.0
ssa	8	1.0	3.1	1.1
tightness0.9	38	1.0	28.1	1.0
varDimacs	16	1.0	3.4	1.1

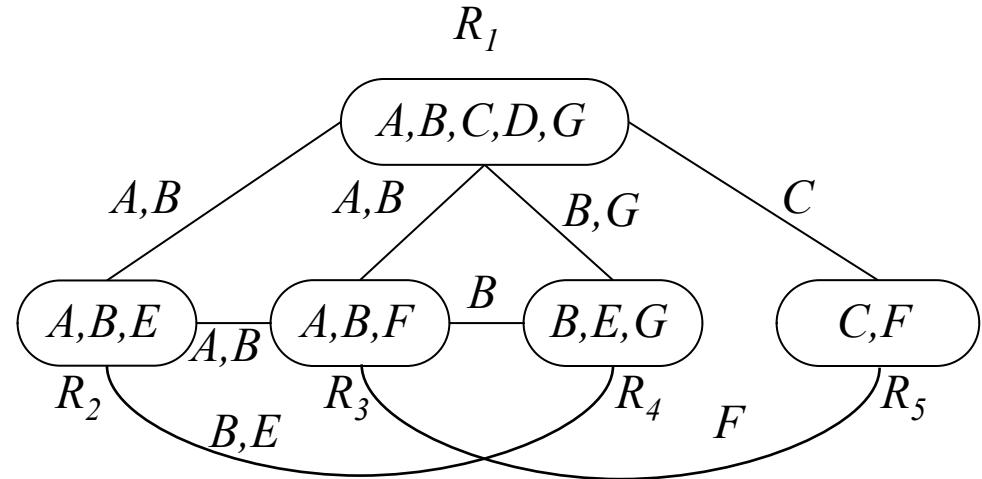
Block Statistics

Benchmark	Abs	Averages		
	Max	Min	Max	Mean
aim-50,100,200,pseudo	4	1.0	4.0	2.1
cmpsed-25-1-2,25,40,80	10	1.0	10.0	8.0-8.4
cmpsed-25-10-20	10	1.0	10.0	7.6
cmpsed-75-1-2,25,40,80	10	1.0	10.0	8.3-8.5
dag-rand	108	1.0	91.6	2.9
dubois,pret	2	1.0	2.0	1.5
geom	20	6.4	20.0	15.0
grCol-hos	6	1.0	3.3	3.3
grCol-mug	3	1.0	2.5	2.4
grCol-register-mulsol	48	23.2	23.2	23.2
grCol-sgb-book	12	1.0	7.7	7.5
grCol-sgb-games	8	1.0	6.3	6.1

Benchmark	Abs	Averages		
	Max	Min	Max	Mean
grCol-sgb-queen	17	10.3	10.3	10.3
hanoi	3	1.0	3.0	2.9
lexVg	875	1.0	484.7	3.6
modifiedRenault	48,720	1.0	48,720.0	7.9
rand-10-20-10	1,046	1.0	119.2	1.3
rand-3-20-20-fcd	190	1.0	181.5	12.8
renault	48,720	1.0	48,720.0	7.7
rlfapGr/ScensMod	44,43	1.0	30.0,35.6	18.5,19.4
ssa	31	1.0	14.7	2.1
super-queens	49	15.6	17.6	16.4
tightness0.9	40	1.0	36.3	16.9
varDimacs	512	1.0	115.0	5.6

Partitioning Relations -- Definitions

- Scope
- Subscope
- Combination

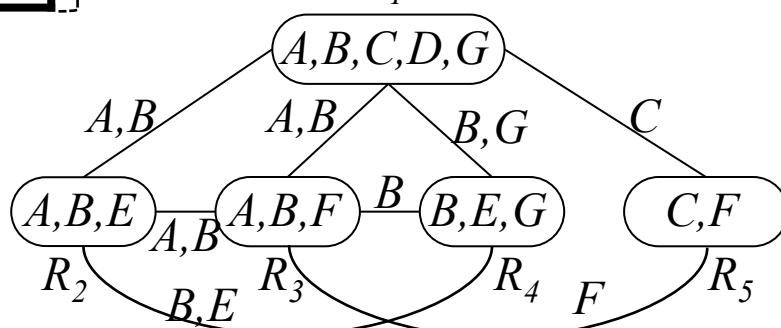


Coarse Partitions

R_1	A	B	C	D	G
fb_1	t_1	0	0	0	0
	t_2	0	0	0	1
fb_2	t_3	0	0	1	0
fb_3	t_4	0	0	1	1
fb_4	t_5	0	1	1	0
	t_5	0	1	1	1
fb_5	t_7	1	1	1	1

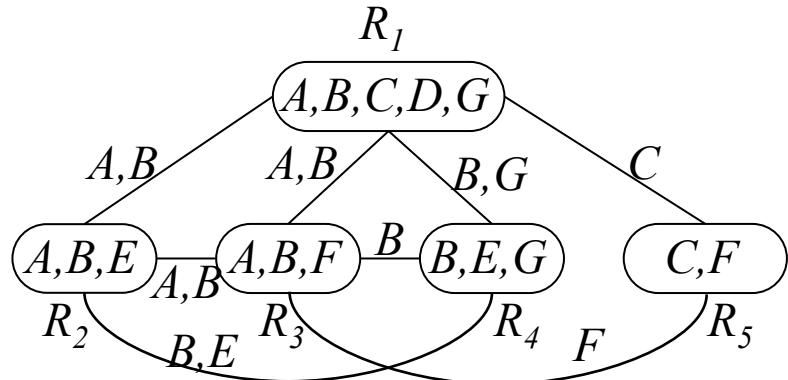
R_2	A	B	E
fb_6	0	0	0
fb_7	0	0	1
fb_8	0	1	0
fb_9	0	1	1
fb_{10}	1	0	0
fb_{11}	1	0	1

R_3	A	B	F
fb_{12}	0	0	0
fb_{13}	0	0	1
fb_{14}	0	1	1
fb_{15}	1	1	0
fb_{16}	1	1	1



Fine Partitions

- Equivalence class that induces smallest set of tuples



R_1	A	B	C	D	G
t_1	0	0	0	0	0
t_2	0	0	0	1	0
t_3	0	0	1	0	0
t_4	0	0	1	1	1
t_5	0	1	1	0	1
t_5	0	1	1	1	1
t_7	1	1	1	1	1

R_2	A	B	E
fb_6	0	0	0
fb_7	0	0	1
fb_8	0	1	0
fb_9	0	1	1
fb_{10}	1	0	0
fb_{11}	1	0	1

Intermediate Partitions

- Consider a combination of size $m=3 \dots$

R_1	A	B	C	D	G	R_1	A	B	E	R_5	C	F
fb_1	t_1	0	0	0	0	0	t_6	0	0	0	0	0
fb_1	t_2	0	0	0	1	0	t_6	0	0	1	0	1
fb_2	t_3	0	0	1	0	0	t_7	0	0	0	0	0
fb_3	t_4	0	0	1	0	1	t_8	0	1	0	0	0
fb_4	t_5	0	1	1	0	1	t_9	0	1	0	0	0
fb_4	t_5	0	1	1	1	1	t_{10}	1	0	0	0	0
fb_5	t_7	1	1	1	1	1	t_{11}	1	0	1	0	0

Intermediate Partitions

- First, identify the subscopes that affect R_1

R_1		A	B	C	D	G
fb_1	t_1	0	0	0	0	0
fb_1	t_2	0	0	0	1	0
fb_2	t_3	0	0	1	0	0
fb_3	t_4	0	0	1	1	1
fb_4	t_5	0	1	1	0	1
fb_4	t_5	0	1	1	1	1
fb_5	t_7	1	1	1	1	1

R_2		A	B	E
fb_6		0	0	0
fb_7		0	0	1
fb_8		0	1	0
fb_9		0	1	1
fb_{10}		1	0	0
fb_{11}		1	0	1

R_5		C	F
fb_{20}		0	0
fb_{21}		0	1
fb_{22}		1	0

Intermediate Partitions

- Project the union of those subscopes over the relation

		R_1						
				A	B	C	D	G
ib_1	fb_1	t_1	0	0	0	0	0	0
		t_2	0	0	0	1	0	
ib_2	fb_2	t_3	0	0	1	0	0	
		t_4	0	0	1	1	1	
ib_3	fb_4	t_5	0	1	1	0	1	
		t_5	0	1	1	1	1	
ib_4	fb_5	t_7	1	1	1	1	1	
		R_2		A	B	E		
				0	0	0		
		R_5		C	F			
				0	0			
				0	0	1		
				0	1	0		
				0	1	1		
				1	0	0		
				1	0	1		